

## Model theory

### 4. Cartesian and reduced products (correction)

#### Exercise 1

1. One can take the language  $\{+, \times, 1, a, b, c\}$ , the structure  $\mathbf{N}$  (with  $+$ ,  $\times$  and  $1$  having their natural interpretation in  $\mathbf{N}$  and the constant symbols  $a, b$  and  $c$  interpreted by  $m, n$  and  $p$ ) and consider the sentence

$$\left( \forall x \forall y (xy = p \rightarrow (x = 1 \vee y = 1)) \right) \wedge \exists u \exists v (um + vn = 1).$$

2. One can take the language of rings  $\{+, \times, -, 0, 1, \leq\}$  augmented with a binary relation symbol, the natural  $L_{ring}$ -structure on  $K$  and  $\leq^K$  the given linear ordering on  $K$  and the sentence

$$\forall x \forall y \forall z (x \leq y \rightarrow x + y \leq x + z) \wedge ((x \geq 0 \wedge y \geq 0) \rightarrow xy \geq 0).$$

3. One can take the language of rings  $\{+, \times, -, 0, 1, \leq\}$  augmented with a binary relation symbol, the natural  $L_{ring}$ -structure on  $K$  and  $\leq^K$  the given linear ordering on  $K$  and the sentence

$$\forall x (x \geq 0 \rightarrow \exists y (x = y^2)).$$

4. One can take the language of rings  $\{+, \times, -, 0, 1, \leq\}$  augmented with 4 constant symbols  $a, b, c, d$ , the natural  $L_{ring}$ -structure on the field  $\mathbf{R}$  of real numbers (with  $a_{11}, a_{12}, a_{21}, a_{22}$  interpreting  $a, b, c, d$ ) and the sentence

$$\exists x \exists y \exists s \exists t (ax + bz = 1 \wedge cy + dt = 1 \wedge ay + bt = 0 \wedge cx + dz = 0),$$

or even the quantifier-free sentence

$$ad - bc = 0.$$

5. One can take the language of groups, the natural  $L_{gp}$ -structure on  $G$  and the sentence

$$\exists x (\forall y (yx = xy) \wedge x \neq 1).$$

#### Exercise 2

Let  $(M_i)_{i \in I}$  be  $L$ -structures  $\varphi(x_1, \dots, x_n)$  an atomic formula and  $a^1, \dots, a^n$  elements of  $\prod_i M_i$ .

##### Claim 2.1

$$\prod_i M_i \models \varphi(a^1, \dots, a^n) \iff M_i \models \varphi(a_i^1, \dots, a_i^n) \text{ for all } i \in I$$

*Proof.* One can show first that, if  $t(\bar{x})$  is an  $L$ -term, writing  $M$  the Cartesian product  $\prod_i M_i$ , then for all  $a^1, \dots, a^n$  in  $M$ , one has  $t^M(a^1, \dots, a^n) = (t^{M_i}(a_i^1, \dots, a_i^n))_{i \in I}$  (by induction on the complexity  $c(t)$  applying the definition of the interpretation of a constant symbol  $c$  and function symbol in  $\prod_i M_i$ ). The Claim follows by applying the definition of the interpretation of a relation symbol  $r$  in  $\prod_i M_i$ .  $\square$

**Claim 2.2** Claim 2.1 does not hold for any formula.

*Proof.* Consider for instance the language  $L$  with equality only, a set  $A$  with two distinct elements, the Cartesian product  $A \times A$  and  $\sigma$  the sentence  $\exists x \exists y \forall z ((x \neq y) \wedge (z = x \vee z = y))$  stating that there are exactly two elements.  $\square$

**Claim 2.3** For any  $J \subset I$ , the restriction map  $\alpha : \prod_{i \in I} M_i \longrightarrow \prod_{j \in J} M_j$  is an  $L$ -morphism.

*Proof.* Let us write  $M_I$  for  $\prod_{i \in I} M_i$  and  $M_J$  for  $\prod_{i \in J} M_i$ . Let  $c$ ,  $r$  and  $f$  be a constant symbol, an  $n$ -ary relation symbol and an  $n$ -ary function symbol respectively. One has

$$\alpha(c^{M_I}) = \alpha((c^{M_i})_{i \in I}) = (c^{M_i})_{i \in J} = c^{M_J},$$

$$\alpha(f^{M_I}(a^1, \dots, a^n)) = \alpha((f^{M_i}(a_i^1, \dots, a_i^n))_{i \in I}) = (f^{M_i}(a_i^1, \dots, a_i^n))_{i \in J} = f^{M_J}(\alpha(a^1), \dots, \alpha(a^n)), \quad \text{and}$$

$$\begin{aligned} (a^1, \dots, a^n) \in r^{M_I} &\iff (\forall i \in I) (a_i^1, \dots, a_i^n) \in r^{M_i} \\ &\implies (\forall i \in J) (a_i^1, \dots, a_i^n) \in r^{M_i} \\ &\iff ((a_i^1)_{i \in J}, \dots, (a_i^n)_{i \in J}) \in r^{M_J} \\ &\iff (\sigma(a^1), \dots, \sigma(a^n)) \in r^{M_J}. \end{aligned} \quad \square$$

Note that Claim 2.1 actually holds for every *positive* formula (i.e. that does not use the negation symbol), using the Axiom of Choice.

### Exercise 3

Let  $I$  and  $J \subset I$  be infinite sets.

**Claim 3.1** There is a non-principal ultrafilter on  $I$  that contains  $\{J\}$ .

*Proof.* Let  $\mathcal{F}$  be the Fréchet filter on  $I$ . Any finitely many elements of  $\mathcal{F} \cup \{J\}$  have a non-empty intersection, so  $\mathcal{F} \cup \{J\}$  generates a filter on  $I$ , which can be extended to an ultrafilter  $\mathcal{U}$  on  $I$ . As  $\mathcal{U}$  contains the Fréchet filter,  $\mathcal{U}$  is non-principal.  $\square$

**Claim 3.2** There is a non-principal ultrafilter on  $\mathbf{N}$  containing  $\{n\mathbf{N} : n \geq 1\}$ .

*Proof.* Let  $\mathcal{F}$  be the Fréchet filter on  $\mathbf{N}$ . Any finitely many elements of  $\mathcal{F} \cup \{n\mathbf{N} : n \geq 1\}$  have a non-empty intersection, so  $\mathcal{F} \cup \{n\mathbf{N} : n \geq 1\}$  generates a filter on  $\mathbf{N}$ , which can be extended to an ultrafilter  $\mathcal{U}$  on  $\mathbf{N}$ . As  $\mathcal{U}$  contains the Fréchet filter,  $\mathcal{U}$  is non-principal.  $\square$

**Claim 3.3** Let  $\mathcal{G}$  be a set of subsets of  $I$ . There is a non-principal ultrafilter extending  $\mathcal{G}$  if and only if  $G_1 \cap \dots \cap G_n$  is infinite for every  $G_1, \dots, G_n$  in  $\mathcal{G}$ .

*Proof.* If  $G_1 \cap \dots \cap G_n$  is infinite for every  $G_1, \dots, G_n$  in  $\mathcal{G}$ , then the above argument holds:  $\mathcal{G} \cup \mathcal{F}$  can be extended to an ultrafilter. Conversely, any non-principal ultrafilter  $\mathcal{U}$  containing  $\mathcal{G}$  must contain the Fréchet filter  $\mathcal{F}$ , and if  $G_1 \cap \dots \cap G_n$  was finite for some  $G_1, \dots, G_n$  in  $\mathcal{G}$ , one would have  $I \setminus (G_1 \cap \dots \cap G_n) \in \mathcal{F}$  so  $(I \setminus (G_1 \cap \dots \cap G_n)) \cap G_1 \cap \dots \cap G_n = \emptyset \in \mathcal{U}$ , a contradiction.  $\square$

## Exercise 4

Let  $I$  be a set,  $J \subset I$  a subset and  $\mathcal{F}$  the principal filter on  $I$  generated by the singleton  $\{J\}$ . Let  $(M_i)_{i \in I}$  be a family of  $L$ -structures.

**Claim 4.1** The reduced product  $\prod_{\mathcal{F}} M_i$  is isomorphic to the Cartesian product  $\prod_{j \in J} M_j$ .

*Proof.* Let  $\alpha : \prod_{\mathcal{F}} M_i \rightarrow \prod_{j \in J} M_j$  be the map sending  $((a_i)_{i \in I})_{\mathcal{F}}$  to  $(a_j)_{j \in J}$ . We claim that  $\alpha$  is well-defined and an  $L$ -isomorphism. If  $((a_i)_{i \in I})_{\mathcal{F}} = ((b_i)_{i \in I})_{\mathcal{F}}$ , then the set  $\{i \in I : a_i = b_i\}$  contains  $J$ , so  $(a_i)_{i \in J} = (b_i)_{i \in J}$  and  $\alpha$  is **well-defined**. If  $(a_i)_{i \in J}$  is an element of  $\prod_{j \in J} M_j$ , let  $(a_i)_{i \in I}$  in  $\prod_{i \in I} M_i$  where  $a_i$  is arbitrarily chosen in  $M_i$  for  $i \in I \setminus J$  (using the axiom of choice). Then  $\alpha$  maps  $((a_i)_{i \in I})_{\mathcal{F}}$  to  $(a_j)_{j \in J}$ , so  $\alpha$  is **surjective**. Let us show that  $\alpha$  is an **embedding**. Let  $c$ ,  $f$  and  $r$  be a constant symbol, an  $n$ -ary function symbol and an  $n$ -ary relation symbol. We write  $M_{\mathcal{F}}$  for  $\prod_{\mathcal{F}} M_i$  and  $M_J$  for  $\prod_{j \in J} M_j$ . For every  $n$ -tuple  $(a^1, \dots, a^n)$  in  $\prod_{i \in I} M_i$ , one has

$$\alpha(c^{M_{\mathcal{F}}}) = ((c^{M_i})_{i \in I})_{\mathcal{F}} = (c^{M_i})_{i \in J} = c^{M_J},$$

$$\alpha(f^{M_{\mathcal{F}}}(a_{\mathcal{F}}^1, \dots, a_{\mathcal{F}}^n)) = \alpha(((f^{M_i}(\bar{a}_i))_{i \in I})_{\mathcal{F}}) = (f^{M_j}(a_j^1, \dots, a_j^n))_{j \in J} = f^{M_J}(\alpha(a_{\mathcal{F}}^1), \dots, \alpha(a_{\mathcal{F}}^n)), \text{ and}$$

$$\begin{aligned} (a_{\mathcal{F}}^1, \dots, a_{\mathcal{F}}^n) \in r^{M_{\mathcal{F}}} &\iff \{i \in I : (a_i^1, \dots, a_i^n) \in r^{M_i}\} \in \mathcal{F} \\ &\iff J \subset \{i \in I : (a_i^1, \dots, a_i^n) \in r^{M_i}\} \\ &\iff (\forall j \in J) (a_j^1, \dots, a_j^n) \in r^{M_j} \\ &\iff (\alpha(a_{\mathcal{F}}^1), \dots, \alpha(a_{\mathcal{F}}^n)) \in r^{M_J}. \end{aligned} \quad \square$$